

Book-keeping --> accounting --> balance --> state

Bookkeeping is the recording of financial transactions, and is part of the process of accounting in business.^[1] Transactions include purchases, sales, receipts and payments by an individual person or an organization/corporation. There are several standard methods of bookkeeping, including the single-entry and double-entry bookkeeping systems.

From <<https://en.wikipedia.org/wiki/Bookkeeping>>

<https://www.dreamstime.com/stock-image-d-life-cycle-accounting-process-illustration-circular-flow-chart-image30625511>



Authorized capital

Credit

Fixed Assets

Costs

Incomes

Op.No. Input Output RemainingAmount

1	123	0	123
2	5	11	117

Compare with UOTX system

<https://medium.com/@olxc/ethereum-and-smart-contracts-basics-e5c84838b19>

State 1

Authorized Capital	Credit	Fixed Assets	Electricity Cost	Mining	Percent for Credit	Balance

State 2

Authorized Capital	Credit	Fixed Assets	Electricity Cost	Mining	Percent for Credit	Balance

State 3

Authorized Capital	Credit	Fixed Assets	Electricity Cost	Mining	Percent for Credit	Balance

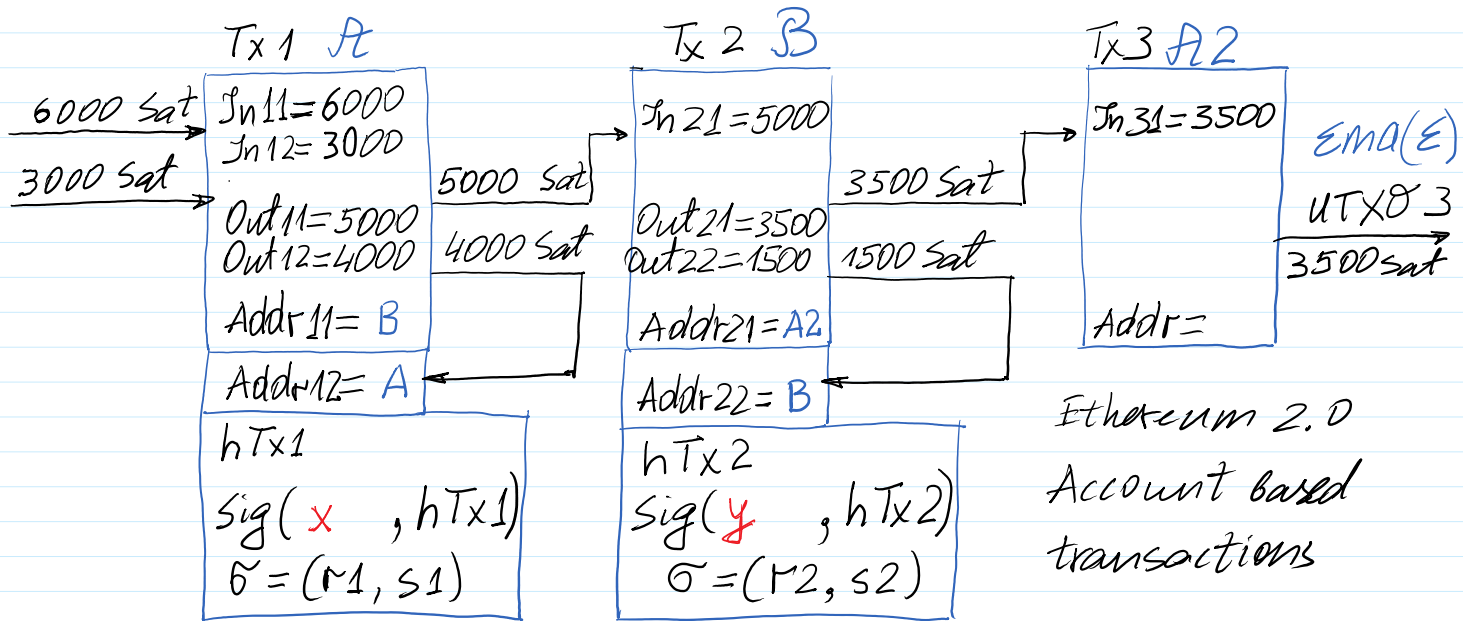
Book-keeping --> Accounting --> Balance --> State

UTxO

Tx 1 A

Tx 2 B

Tx 3 A2



'Tx1 : $In11 = 6000 \parallel In12 = 3000 \parallel Out11 = 5000 \parallel Out12 = 4000 \parallel Rec1 = B \parallel Rec2 = A$ '

'Tx2 : $In21 = 5000 \parallel Out21 = 3500 \parallel Out22 = 1500 \parallel Rec1 = A2 \parallel Rec2 = B$ '

'Tx3 : $In31 = 3500 \parallel Out31 = 3500 \parallel Rec = E$ '

$Tx1 = '---'$ $Tx2 = '---'$ $Tx3 = '---'$

$h_1 = H(Tx1) = h28(Tx1)$

$h_2 = H(Tx2) = h28(Tx2)$

$h_3 = H(Tx3) = h28(Tx3)$

Transaction template:

$Tx_N = 'TxN:In11=... \parallel In12=... \parallel Out11=... \parallel Out12=... \parallel Rec1=... \parallel Rec2=...'$

Transactions:

$Tx_1 = 'Tx1:In11=6000 \parallel In12=3000 \parallel Out11=5000 \parallel Out12=4000 \parallel Rec1=B \parallel Rec2=A'$

$Tx_2 = 'Tx2:In21=5000 \parallel Out21=3500 \parallel Out22=1500 \parallel Rec1=A2 \parallel Rec2=B'$

$Tx_3 = 'Tx3:In31=3500 \parallel Out31=3500 \parallel Out32=0 \parallel Rec1=E \parallel Rec2=A2'$

$>> hTx_1 = h28('Tx1:In11=6000 \parallel In12=3000 \parallel Out11=5000 \parallel Out12=4000 \parallel Rec1=B \parallel Rec2=A')$

$hTx_1 = 5B5412B$

$>> hTx_1 = h28(Tx_1)$

$hTx_1 = 5B5412B$

>> hTx_2=h28('Tx2:ln21=5000||Out21=3500||Out22=1500||Rec1=A2||Rec2=B')

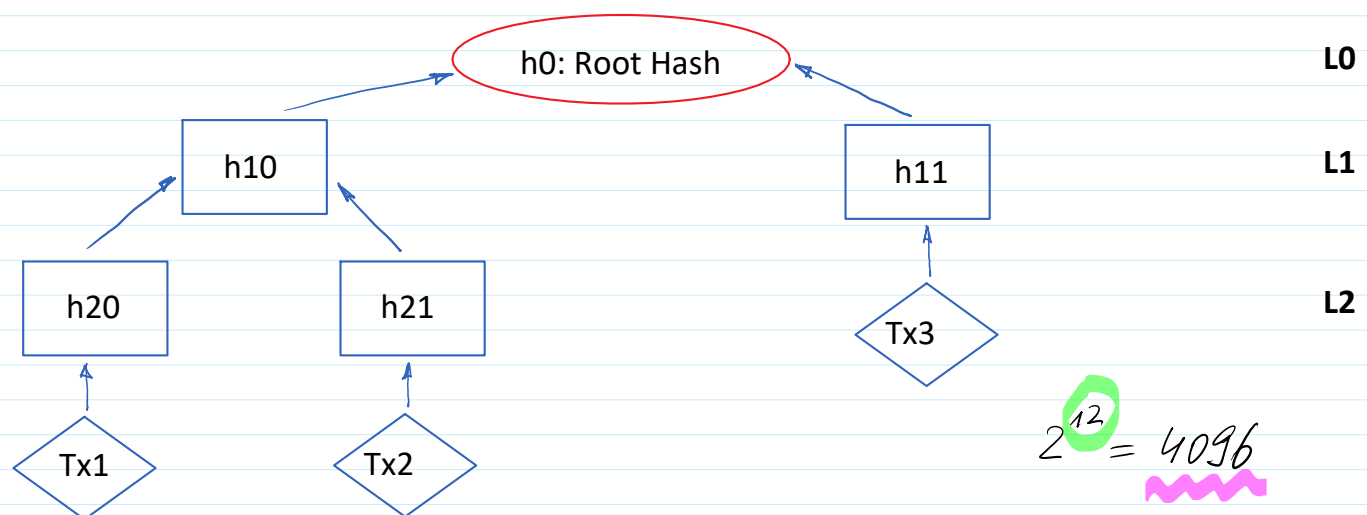
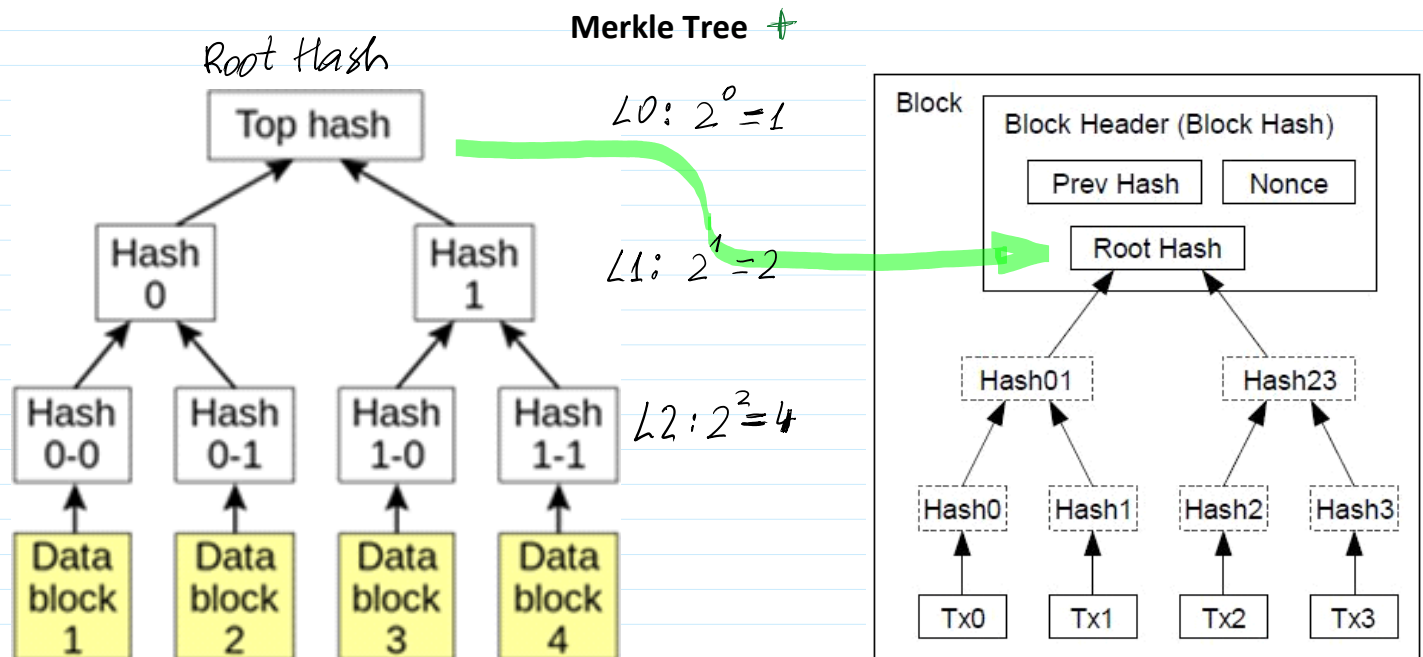
>> hTx_2=h28(Tx_2)

hTx_2 = D5C895A

>> hTx_3=h28('Tx3:ln31=3500||Out31=3500||Out32=0||Rec1=E||Rec2=A2')

>> hTx_3=h28(Tx_3)

hTx_3 = FEC59B7



>> h20=h28(hTx_1)

h20 = 5B5412B

>> h21=h28(hTx_2)

h21 = D5C895A

>> h10=h28('5B5412B||D5C895A')

h10 = B6BD3A4

Root Hash: h0

```

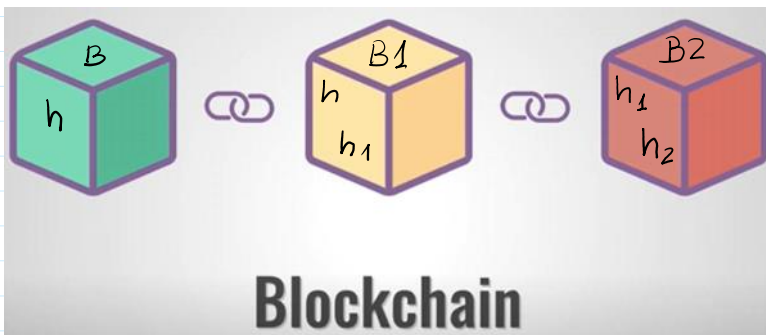
h20 = 5B5412B
>> h21=h28(hTx_2)
h21 = D5C895A
>> h11=h28(hTx_3)
h11 = FEC59B7
>> h10=h28('5B5412B||D5C895A')
h10 = B6BD3A4
Root Hash: h0
>> h0=h28('B6BD3A4||FEC59B7')
h0 = 60BA3B5

```

Python : sha256

h20: 5B5412B
h21: D5C895A
h11: FEC59B7

h10: 625A41F
h0: **60BA3B5**



Magic Number (4)		Block Size (4)	
Version (4)		Previous Block Hash (32)	
		Merkle Root(32)	
		Timestamp (4)	
Difficulty Target (4)		Nonce (4)	
Transaction Counter (Variable : 1-9)			
Transaction List (Variable : Upto 1 MB)			

BLOCK HEADER

Block size = 4 Bytes
4 Bytes x 8 bits = 32 bits
Block have
 $2^{32} - 1 = 4294967295$
In ASCII encoding
8 bits represents
1 symbol a, b, c, --#.
Block represents
536 870 912 symbols

Difficulty Target (DT): defines the complexity of block mining. In our simulation DT we will choose to find h-value of mining (mined block) having only 1 leading hexadecimal digit equal to 0.

$h28('RootHash_PrevHash_737327631') =$

>> sha256('RootHash PrevHash 737327631')

ans = F4AE534CD226FAF7998C8424B348E020BA80639A687E93A0B8C5130ED C51E6DE

C51E6DE

>> sha256('RootHash PrevHash 737327632')

ans = B856211DF2EE15E30AB770C1A43CE014ECFE573182AFD885B28D96854DBC5F21

>> sha256('RootHash PrevHash 737327633')

ans = 9C18C764E347A58E57AC3F7A3C2874D5889A0E802699FEA47EEFF8C03BFEDA69

DT: to mine a block it is needed to find h-value having leading zero in hexadecimal format: C51E6DE OXXXXXX

6x4=24 bits

h-value is computed so >> h28() → 7 hex numbers

What probability to mine a block?

The number of possible h-values of 28 bits length

>> 2^{28} ans = 268 435 456

The number of adequate h-values: 2^{24}

$$Pr\{to Mine\} = \frac{2^{24}}{2^{28}} = \frac{1}{2^4} = \frac{1}{16}$$

DT: two leading hex number = 00

The number of adequate h-values: 2^{20}

00XXXXX

5x4=20

$$Pr\{to Mine\} = \frac{2^{20}}{2^{28}} = \frac{1}{2^8} = \frac{1}{256}$$

$$Pr\{to Mine\} = \frac{2^{12}}{2^{28}} = \frac{1}{2^{16}} = \frac{1}{65536}$$

>> 2^{16} ans = 65536

$$Pr\{to Mine\} = \frac{1}{2^{28}} = \frac{1}{268\,435\,456}$$

>> 2^{28} ans = 268 435 456

The probability to mine a block, e.g. in Bitcoin when

DT: is to find SHA256 value having 18 leading zeroes

```
>> sha256('RootHash PrevHash 737327631')
```

```
ans = F4AE534CD226FAF7998C8424B348E020BA80639A687E93A0B8C5130EDC51E6DE  
00000000000000000000 XXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXX
```

The number of possible h -values having 256 bits is 2^{256} .

The number of adequate h -values of SHA 256 is

$256 - 18.4 = 256 - 72 = 184$ bits, that are represented 46 hex. num.

The number of adequate values is 2^{184} .

$$\text{Prob}\{\text{to mine}\} = \frac{2^{184}}{2^{256}} = 2^{184-256} = 2^{-72}$$

$$2^{72} \sim 4 \text{ GT} = 4 \cdot 2^{30} \cdot 2^{40} = 2^2 \cdot 2^{30} \cdot 2^{40} = 2^{72}$$

$$1 \text{ K} = 2^{10} = 1024$$

$$1 \text{ M} = 2^{20} = \dots$$

$$1 \text{ G} = 2^{30} = \dots$$

$$1 \text{ T} = 2^{40} = \dots$$