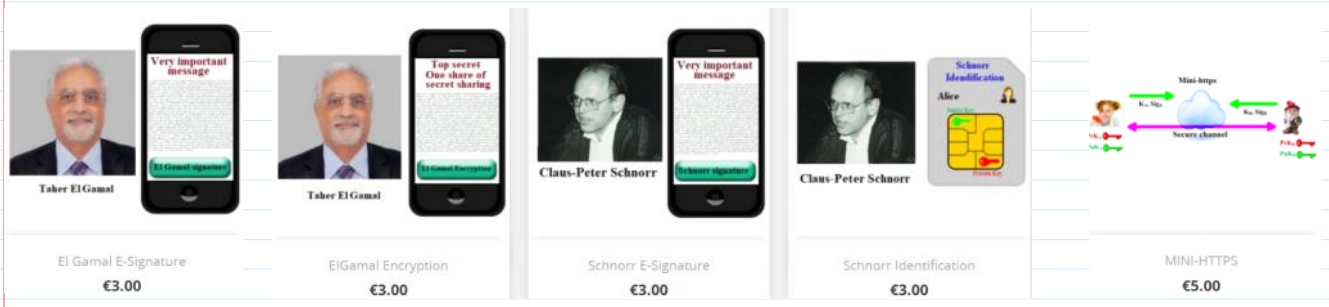
For Poster Report (PR) presentation 2 lectures will be dedicated at the end of semester.
During this time you can pass the Mid-Term Exam (MTE) if it is not passed yet.

PR requirements you can find in my Google drive:
https://docs.google.com/document/d/1IPzwEVVmvnObQoJPRP9GDUiHccEGWgMw/edit?usp=sharing&ouid=111502255533491874828&rtpof=true&sd=true

PR Topics are presented in my Google drive:
https://docs.google.com/document/d/1B6gavCsgZXcCRssFZEWLVfzaO_IPbC5o/edit?usp=sharing&ouid=111502255533491874828&rtpof=true&sd=true

Please select a topic you wish and label this selection in the list.

During the exam you must solve 5 problems from
https://imimsociety.net/en/14-cryptography
And oral answer to the 1-2 questions concerning these 5 problems.




**Information Confidentiality, Integrity and Authenticity. Person Identification.**

**Public Parameters PP = ($p$, $g$).**
**p = 268435019; g=2;**

**Information**
**Encryption/Decryption; Signing/Verification**
**Person Identification using Zero Knowledge Proof - ZKP**

**Interactive Zero Knowledge Proof - ZKP**

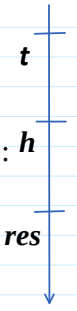
**$\mathcal{A}$: ZKP of knowledge $x$:**
$PrK_A = x$ = randi($p$-1)
$PuK_A = a = g^x \bmod p$
1.Computes **commitment** $t$ for random number $i$:
$i$=randi($p$-1)
$t=g^i \bmod p$
3.Computes **response** $res$:
$res=i+xh \bmod (p-1)$

$t, a$ →
← $h$
$res$ →

$\mathcal{B}$: $PuK_A = a$

2.Generates **challenge** $h$: $h$
$h$=randi($p$-1)

Verifies:
$g^{res}=ta^h \bmod p$



**Alice**: 'Hello Bob'
>> M='Hello Bob'
***M***; **σ=(*r*,*s*)**; ***a***

***h*=H(*M*||*r*)**.

***M'***; **σ=(*r*,*s*)**; ***a***

**Bob:** let ***M'*=*M*.**
1.Computes ***h*=H(*M*||*r*)**.
>> h=concat(M,r)
2.Verifies signature on ***h***.

**Sign(PrK=*a*, *PP*, *h*) = σ = (*r*, *s*)**

**Ver(*a*,σ,*h*)=V∈{*True*, *False*}≡{1, 0}.**

**Non-Interactive Zero Knowledge Proof - NIZKP.**

**Alice** using the Schnorr Signature scheme can prove a knowledge of any other secret in one or other way related with Discrete Exponent Function - DEF.
Let this secret is some integer ***i*** and then **Alice** using DEF computes so called **Statement** we denote by ***t*** for her secret ***i***:

$$t = g^i \bmod p.$$

In this scenario Alice is called a ***Prover***.
Then **Alice** realizes a NIZKP of knowledge of ***i*** without revealing ***i*** by presenting this **Statement** ***t*** to the ***Verifier* Bob**.

A: $u \leftarrow randi(\mathbb{Z}_{p-1})$; $\mathbb{Z}_{p-1} = \{0, 1, 2, \dots, p-2\}$; $+, -, *, :\ mod(p-1)$

$r = g^u \bmod p$
$h = H(a||t||r)$
$s = u + ih \bmod (p-1)$
$\Pi = (r, s)$

$\Pi=(r,s), t, a$ →

B: $h = H(a||t||r)$
$g^s = r \cdot t^h \bmod p$

$$g^s = g^{u+ih \bmod (p-1)} \bmod p = g^u \cdot g^{ih} \bmod p = r \cdot (g^i)^h = r \cdot t^h \bmod p$$

**Mini-https**

**PrK$_A$ = *x* = randi(*p*-1)**
**PuK$_A$ = *a* = $g^x$ mod *p***

Hand shaking $PuK_A$ ⟷
AKAP ⟷
$G$, $\sigma=(r,s)$ ⟷

**PrK$_B$ = *y* = randi(*p*-1)**
**PuK$_B$ = *b* = $g^y$ mod *p***

**PuK$_A$ = *a***

$u \leftarrow randi(p-1)$
>> $u$ = int64(randi(p-1))
$t_A = g^u \bmod p$
>> $t_A$ = mod_exp(g, u, p)
Sign($x$, $t_A$) = $\sigma_A = (r_A, s_A)$

$t_A$, $\sigma_A = (r_A, s_A)$ →
$PuK_A$

1. Verify if $PuK_A$ is in Data Base.
2. Verify if $\sigma_A$ on $t_A$ is valid.
Ver($PuK_A$, $\sigma_A$, $t_A$) = T

$Sign(x, t_A) = \sigma_A = (r_A, s_A)$ $\quad PuK_A$

$i \leftarrow randi(p-1)$

$r_A = g^i \bmod p$

2. Verify if $\sigma_A$ on $t_A$ is valid.

$Ver(PuK_A, \sigma_A, t_A) = T$

$t_B, \sigma_B$ / $PuK_B$

3. Generates $v \leftarrow int64(randi(p-1))$

computes $t_B = g^v \bmod p$.

$sign(y, t_B) = \sigma_B = (r_B\ s_B)$

$\sigma_B = (R, S)$

Ver(PuK$_B$ = $b$, $t_B$) = True

$k_{AB} = (t_B)^u \bmod p =$

$= (g^v)^u \bmod p = g^{vu} \bmod p =$

$k_{AB} = k = k_{BA}$

$k_{BA} = (t_A)^v \bmod p =$

$= (g^u)^v \bmod p = g^{uv} \bmod p$

A: creates transaction $T_x$

$E(k, T_x) = C$ $\quad C, \sigma = (r, s)$

$j \leftarrow$ int64(randi($p$-1))

r = $g^j$ mod $p$

$h = H(C \| r)$

$s = j + x\,h \bmod (p-1)$

$\sigma = (r, s)$

1. $Ver(PuK_A = a, \sigma, C) = \{T, F\}$

2. $D(k, C) = T_x$

3. Performes money transf.


MINI-HTTPS

€5.00

After receiving ***Tx*** and **σ**, **Bob** according computes ***h***

***h* = H(*C*||*r*)**,

and verifies if

$g^s \bmod p = r a^h \bmod p$.

V1 V2

Symbolically this verification function we denote by

**Ver(*a*,σ,*h*)=V∈{*True*, *False*}≡{1, 0}**. This function yields ***True*** if (2.22) is valid if: **PuK$_A$= *a* =*F*(PrK$_A$)= $g^x$ *mod p***

1. Mentor sends you Public Parameters (***p***=268435019; ***g***=2) of 28 bits length.
Generate and private public keys **PrK$_A$=*x*** and **PuK$_A$=*a***.
Send public key [***a***] to the Mentor.

```
>> a
a = 39794323
```

39794323

2. Compute random secret number ***u*** of 28 bit length and compute session public parameter $t_A$. Sign $t_A$ with Schnorr signature scheme by computing two signature components $\sigma_A = (r_A, s_A)$. Send [tA,rA,sA] to the Mentor.

```
>> u=int64(randi(p))
u = 190442301
>> tA=mod_exp(g,u,p)
tA = 109819746
>>
>> i=int64(randi(p))
i = 236712983
```

```
% Alice verifies her signature before sending to Mentor
>> g_sA=mod_exp(g,sA,p)
g_sA = 254713335
```

% Alice verifies her signature before sending to Mentor
>> g_sA=mod_exp(g,sA,p)
g_sA = 254713335
>> V1=g_sA
V1 = **254713335**
>> a_h=mod_exp(a,h,p)
a_h = 85497572
>> V2=mod(rA*a_h,p)
V2 = **254713335**

109819746,117664796,114109665

>>
>> i=int64(randi(p))
i = 236712983
>> rA=mod_exp(g,i,p)
r A= 117664796
>> con=concat(tA,r)
con = 109819746117664796
>> h=hd28(con)
h = 243151357
>> sA=mod(i+x*h,p-1)
sA = 114109665

3. Mentor sends you ($t_B$ **PuK$_B$**=32768, $K_B$ $t_B$=209399419, $R$=101644938, $S$=18011748). Verify Mentor's signature $\sigma_M=(R,S)$ on $t_B$. If signature is valid then taking $S$ compute verification parameter $V_1=g^S \bmod p$. Compute common symmetric secret key **k** and transform **k** to the hexadecimal form ***kh*** of 32 digits length as it is required for AES128 function. Create the string of message variable ***m***='MMDD' consisting of the month and day of your birth. Encrypt message ***m*** using 1 round of AES128 cipher with key ***kh***32 by computing ciphertext >> C=AES128(m,kh32,1,'e'). **Attention**! Encryption using 1 round is extremely insecure and is used there to speed up the computations and to make sure of its insecurity. Insecurity is seen by comparing plaintext and ciphertext messages in hexadecimal format. They have non-excrypted digits. ***C*** should be entered within ' '. Send [V1,C] to the Mentor for decryption.

>> PuKB=int64(32768)
PuKB = 32768
>> tB=int64(209399419)
tB = 209399419
>> R=int64(101644938)
R = 101644938
>> S=int64(18011748)
S = 18011748
>>
>> con=concat(tB,R)
con = 209399419101644938
>> h=hd28(con)
h = 64128805

% Alice verifies Bank's signature
>> g_S=mod_exp(g,S,p)
g_S = 20703551
>> V1=g_S
V1 = 20703551
V2 = 20703551

% Alice computes common
% symmetric secret key ***k***
>> k=mod_exp(tB,u,p)
k = 63198998

% AES128(in,kh32,NR,fun) Advanced Encryption Standard symmetric cipher with key length of 128 bits
% Encryption is performed for 1 block of length 128 bits or 16 ASCII symbols
%
% in - plaintext/ciphertext of string type: maximum 16 symbols or shorter
%
% kh32 - shared secret key in hexadecimal number of length=32 (128 bits)
% kh32 can be obtained when shared decimal key k is given using commands:
% >> k=int64(randi(2^28))
% k = 160966896
% >> kh32=dec2hex(k,32)
% kh32 = 00000000000000000000000099828F0
%
% NR - Number of Rounds (e.g. Nr = 10)
% The smaller NR, the lower security of encryption but the speed of encryption is higher
% The least number of NR is 1 and in this case security lack is evident
%
% fun - letter determining either encription: fun='e' or decryption: fun='d' functions

>> kh32=dec2hex(k,32)
kh32 = 0000000000000000000000003C45716
>> m='1012'
m = 1012
>> NR=1
NR = 1
>> C=AES128(m,kh32,NR,'e')
new = ~8$M~8t ~[illegible] =[illegible]
C = 7e38244d7e3874187ee424183dfc730e

20703551,'7e38244d7e3874187ee424183dfc730e'

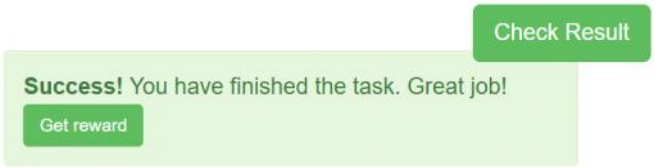
4. Ok, let be informed that Mentor gets you a price for your birthday. The sum of the price he is sending to you as a ciphertex

>> CM='7e38245d7e38d8187e47241865fc730e'
CM = 7e38245d7e38d8187e47241865fc730e

4. Ok, let be informed that Mentor gets you a price for your birthday. The sum of the price he is sending to you as a ciphertex (***CM***='7e38245d7e38d8187e47241865fc730e'). Please decrypt and check it and then encrypt it again with added string 'ok' right after the sum by computing ciphertext ***C***1.
Send [C1] to the Mentor. ***C***1 should be entered within ' '.

'7e3824847e3840187efa24189efc730e'

```
>> CM='7e38245d7e38d8187e47241865fc730e'
CM = 7e38245d7e38d8187e47241865fc730e
>>
>> M=AES128(CM,kh32,NR,'d')
Out = 00000000000000000000000000003935
M = 95
>>
>> Mok='95ok'
Mok = 95ok
>>  C1=AES128(Mok,kh32,NR,'e')
new = ~8$8@  ~
C1 = 7e3824847e3840187efa24189efc730e
```




Till this place